

Bergische Universität Wuppertal

Fachbereich Mathematik und Naturwissenschaften

Institute of Mathematical Modelling, Analysis and Computational Mathematics (IMACM)

Preprint BUW-IMACM 19/20

Lorenc Kapllani, Long Teng and Matthias Ehrhardt

A Multistep Scheme to solve Backward Stochastic Differential Equations for Option Pricing on GPUs

June 2019

http://www.math.uni-wuppertal.de

A Multistep Scheme to solve Backward Stochastic Differential Equations for Option Pricing on GPUs

Lorenc Kapllani, Long Teng and Matthias Ehrhardt

Abstract The goal of this work is to parallelize the multistep method for the numerical approximation of the Backward Stochastic Differential Equations (BSDEs) in order to achieve both, a high accuracy and a reduction of the computation time as well. In the multistep scheme the computations at each grid point are independent and this fact motivates us to select massively parallel GPU computing using CUDA. In our investigations we identify performance bottlenecks and apply appropriate optimization techniques for reducing the computation time, using a uniform domain. Finally, a Black-Scholes BSDE example is provided to demonstrate the achieved acceleration on GPUs.

1 Introduction

In this work we parallelize the multistep scheme developed in [19] to approximate numerically the solution of the following (decoupled) *forward backward stochastic differential equation (FBSDE)*:

$$\begin{cases} dX_t &= a(t, X_t) dt + b(t, X_t) dW_t, \quad X_0 = x_0, \\ -dy_t &= f(t, X_t, y_t, z_t) dt - z_t dW_t, \\ y_T &= \xi = g(X_t), \end{cases}$$
(1)

Lorenc Kapllani

Long Teng

University of Wuppertal, Applied Mathematics and Numerical Analysis, Gaussstrasse 20, D-42119 Wuppertal, Germany, e-mail: kapllani@math.uni-wuppertal.de

University of Wuppertal, Applied Mathematics and Numerical Analysis, Gaussstrasse 20, D-42119 Wuppertal, Germany, e-mail: teng@math.uni-wuppertal.de

Matthias Ehrhardt

University of Wuppertal, Applied Mathematics and Numerical Analysis, Gaussstrasse 20, D-42119 Wuppertal, Germany, e-mail: ehrhardt@math.uni-wuppertal.de

where $X_t, a \in \mathbb{R}^n$, *b* is a $n \times d$ matrix, W_t is a *d*-dimensional Brownian motion, $f(t, X_t, y_t, z_t) : [0, T] \times \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^{m \times d} \to \mathbb{R}^m$ is the driver function and ξ is the terminal condition. We see that the terminal condition y_T depends on the final value of a forward stochastic differential equation (SDE). For a = 0 and b = 1, namely $X_t = W_t$, one obtains a *backward stochastic dierential equation (BSDE)* of the form

$$\begin{cases} -dy_t = f(t, y_t, z_t)dt - z_t dW_t, \\ y_T = \xi = g(W_T), \end{cases}$$
(2)

where $y_t \in \mathbb{R}^m$ and $f(t, y_t, z_t) : [0, T] \times \mathbb{R}^m \times \mathbb{R}^{m \times d} \to \mathbb{R}^m$. In the sequel of this work, we investigate the acceleration of numerical scheme developed in [19] for solving (2). Note that the developed schemes can be applied also for solving (1), where the general Markovian diffusion X_t can be approximated, e.g., by using the Euler-Scheme.

The existence and uniqueness of the solution of (2) were proven by Pardoux and Peng [12]. Peng [13] obtained a direct relation between forward-backward stochastic differential equations (FBSDEs) and partial differential equations (PDEs). Based on this relationship, many numerical methods are proposed, e.g. probabilistic based methods in [3, 4, 8, 11, 20], tree-based methods in [5, 16] etc. El Karoui, Peng and Quenez [7] showed that the solution of a linear BSDE is in fact the pricing and hedging strategy of an option derivative. This was the first claim of application of BSDEs in finance.

In general the solution of BSDEs cannot be established in a closed form. Therefore, a numerical method is mandatory. There are two main classes of numerical methods for approximating the solution of BSDEs. The first class is related with the PDE equivalent based on the Feynman-Kac formula and the second is based on the BSDE. Many methods have been developed, but one of the most interesting (due to the ability to achieve very high accuracy) is developed by Zhao, Zhang and Ju [19]. They used Lagrange interpolating polynomials to approximate the integrals, given the values of integrands at multiple time levels. One of the drawbacks of their method is the computation time. However, the method is highly parallel. Hence, due to simple and intense calculations, the best computing environment is the one offered by GPUs rather than CPUs.

Many acceleration strategies have been developed to solve option pricing problems on the GPU with different mathematical models. However, little work is based on BSDEs. Dai, Peng and Dong [6] solved a linear BSDE on the GPU with the thetascheme method. They analyzed the effects of the thread number per block to increase the speedup. The parallelized program using CUDA achieved high speedups and showed that the GPU architecture is well suited for solving the BSDEs in parallel. Later in 2011, they developed acceleration strategies for option pricing with nonlinear BSDEs using a binomial lattice based method [14]. To increase the speedup, they reduce the global memory access frequency by avoiding the kernel invocation on each time step. Also, due to the load imbalance produced by the binomial grid, they provided load-balanced strategies and showed that the acceleration algorithms exhibit very high speedup over the sequential CPU implementation and therefore suitable for real-time application.

In 2014, Peng, Liu, Yang and Gong [15] considered solving high dimensional BSDEs on GPUs with application in high dimensional American option pricing. A *Least Square Monte-Carlo (LSMC)* method based numerical algorithm was studied, and summarised in four phases. Multiple factors which affect the performance (task allocation, data store/access strategies and the thread synchronisation) were taken into consideration. Results showed much better performance than the CPU version. In 2015, Gobet, Salas, Turkedjiev and Vasquez [9] designed a new algorithm for solving BSDEs based on LSMC. Due to stratification, the algorithm is very efficient especially for large scale simulations. They showed big speedups even in high dimensions.

Next we introduce some preliminary elements which are needed to understand the multistep scheme. We start with the relation of BSDEs and PDEs. Let $(\Omega, \mathscr{F}, \mathbb{P}, \{\mathscr{F}_t\}_{0 \le t \le T})$ be a complete, filtered probability space. In this space a standard *d*-dimensional Brownian motion W_t is defined, such that the filtration $\{\mathscr{F}_t\}_{0 \le t \le T}$ is generated. We define $\|\cdot\|$ as the standard Euclidean norm in the Euclidean space \mathbb{R}^m or $\mathbb{R}^{m \times d}$ and $L^2 = L_F^2(0,T;\mathbb{R}^d)$ the set of all $\{\mathscr{F}_t\}$ -adapted and square integrable processes valued in \mathbb{R}^d . A pair of processes $(y_t, z_t) : [0,T] \times \Omega \to$ $\mathbb{R}^m \times \mathbb{R}^{m \times d}$ is the solution of BSDE (2) if it is $\{\mathscr{F}_t\}$ -adapted, square integrable, and satisfies (2) in the sense of

$$y_t = \xi + \int_t^T f(s, y_s, z_s) \, ds - \int_t^T z_s \, dW_s, \quad t \in [0, T), \tag{3}$$

where $f(t, y_t, z_t) : [0, T] \times \mathbb{R}^m \times \mathbb{R}^{m \times d} \to \mathbb{R}^m$ is $\{\mathscr{F}_t\}$ -adapted and the third term on the right-hand side is an Itô-type integral. This solution exist under "reasonable" regularity conditions [12]. Let us consider the following:

$$y_t = u(t, W_t), \quad z_t = \nabla u(t, W_t) \quad \forall t \in [0, T),$$
(4)

where ∇u denotes the derivative of u(t,x) with respect to the spatial variable *x* and u(t,x) is the solution of the following (backward in time) parabolic PDE:

$$\frac{\partial u}{\partial t} + \frac{1}{2} \sum_{i=1}^{d} \frac{\partial^2 u}{\partial x_i^2} + f(t, u, \nabla u) = 0,$$
(5)

with the terminal condition $u(T,x) = \phi(x)$. Under "reasonable" conditions, the PDE (5) possess a unique solution u(t,x). Therefore, for $\xi = \phi(W_T)$, the pair (y_t, z_t) is the unique solution of BSDE (3).

Due to conditional expectations that compound the numerical method, the following notations will be used. Let $\mathscr{F}_s^{t,x}$ for $t \le s \le T$ be a σ -field generated by the Brownian motion $\{x + W_r - W_t, t \le r \le s\}$ starting from the time-space point (t, x). We define $E_s^{t,x}[X]$ as the conditional expectation of the random variable *X* under the filtration $\mathscr{F}_s^{t,x}$, i.e. $E_s^{t,x}[X] = E[X|\mathscr{F}_s^{t,x}]$. This work is organized as follows. In Section 2 we introduce the multistep scheme. Next, in Section 3 our algorithmic framework for using GPU is presented. In Section 4 we illustrate our findings with the Black-Scholes example.

2 The Multistep scheme

In this section we briefly present the multistep scheme. This is done in two steps, the first corresponds to the derivation of the stable semi-discrete scheme, as only the time domain is discretized. Furthermore, the space is discretized and the fully stable multistep scheme is achieved. Note that the scheme will be presented for the one-dimensional case (but recall that in principle it can be generalized for the *d*-dimensional case).

2.1 The stable semi-discrete scheme

Let *N* be a positive integer and $\Delta t = T/N$ the step size that partitions uniformly the time interval [0,T]: $0 = t_0 < t_1 < \cdots < t_{N-1} < t_N = T$, where $t_i = t_0 + i\Delta t$, $i = 0, 1, \dots, N$. Let *k* and K_y be two positive integers such that $1 \le k \le K_y \le N$. The BSDE (3) can be expressed as

$$y_{t_n} = y_{t_{n+k}} + \int_{t_n}^{t_{n+k}} f(s, y_s, z_s) \, ds - \int_{t_n}^{t_{n+k}} z_s \, dW_s.$$
(6)

In order to approximate y_{t_n} based on the later information $[t_n, t_{n+k}]$, we need to adapt it to the filtration (that is already generated, since we are solving it backwards). Therefore, taking the conditional expectation $E_{t_n}^x[\cdot]$ in (6), we have

$$y_{t_n} = E_{t_n}^x \left[y_{t_{n+k}} \right] + \int_{t_n}^{t_{n+k}} E_{t_n}^x \left[f(s, y_s, z_s) \right] ds$$
(7)

where the third term of (6) is disappeared as it is an Itô-type integral. In order to approximate the integral part of (7), Zhao [19] considered the Lagrange interpolating method, since the $E_{t_n}^x [f(s, y_s, z_s)]$ is a deterministic function of *s*. Given the values of $(t_{n+i}, E_{t_n}^x [f(t_{n+i}, y_{t_{n+i}}, z_{t_{n+i}})])$ and using Lagrange interpolating polynomial, (7) becomes

$$y_{t_n} = E_{t_n}^{x} \left[y_{t_{n+k}} \right] + k\Delta t \sum_{i=0}^{K_y} b_{K_y,i}^{k} E_{t_n}^{x} \left[f(t_{n+i}, y_{t_{n+i}}, z_{t_{n+i}}) \right] + R_y^n,$$
(8)

where $b_{K_y,i}^k$ are the coefficients derived from the integration of Lagrange interpolating polynomial [19] and R_y^n is the error due to the former.

4

A Multistep Scheme to solve BSDEs for Option Pricing on GPUs

Next, we derive a semi-discretized form for the z_t process. Let $\Delta W_s = W_s - W_{t_n}$ for $s \ge t_n$. Then ΔW_s is a standard Brownian motion with mean 0 and standard deviation $\sqrt{s-t_n}$. Let l and K_z be two positive integers such that $1 \le l \le K_z \le N$. Using l instead of k in (6), multiplying both sides by $\Delta W_{t_{n+l}}$, taking the conditional expectation $E_{t_n}^{x}[\cdot]$ and using the Itô isometry we obtain

$$0 = E_{t_n}^x \left[y_{t_{n+l}} \Delta W_{t_{n+l}} \right] + \int_{t_n}^{t_{n+l}} E_{t_n}^x \left[f(s, y_s, z_s) \Delta W_s \right] ds - \int_{t_n}^{t_{n+l}} E_{t_n}^x \left[z_s \right] ds.$$
(9)

Using again the Lagrange interpolation method to approximate the two integrals in (9), we have

$$0 = E_{t_n}^x \left[y_{t_{n+l}} \Delta W_{t_{n+l}} \right] + l \Delta t \sum_{i=0}^{K_z} b_{K_z,i}^l E_{t_n}^x \left[f(t_{n+i}, y_{t_{n+i}}, z_{t_{n+i}}) \Delta W_{t_{n+i}} \right] + R_{z_1}^n$$

$$- l \Delta t \sum_{i=0}^{K_z} b_{K_z,i}^l E_{t_n}^x \left[z_{t_{n+i}} \right] - R_{z_2}^n,$$
(10)

where $b_{K_{z,i}}^{l}$ are the coefficients derived from the integration of the Lagrange interpolating polynomial and (R_{z1}^{n}, R_{z2}^{n}) are the errors for the first and second integrals in (9).

Consider (y^n, z^n) as an approximation of (y_t, z_t) , the semi-discrete scheme is defined as follows: Given random variables (y^{N-i}, z^{N-i}) , i = 0, 1, ..., K - 1 with $K = \max\{K_y, K_z\}$, find the random variables (y^n, z^n) , n = N - K, ..., 0 such that

$$y^{n} = E_{t_{n}}^{x} \left[y^{n+k} \right] + k\Delta t \sum_{i=0}^{K_{y}} b_{K_{y},i}^{k} E_{t_{n}}^{x} \left[f(t_{n+i}, y^{n+i}, z^{n+i}) \right]$$

$$0 = E_{t_{n}}^{x} \left[z^{n+l} \right] + \sum_{i=1}^{K_{z}} b_{K_{z},i}^{l} E_{t_{n}}^{x} \left[f(t_{n+i}, y^{n+i}, z^{n+i}) \Delta W_{t_{n+i}} \right] - \sum_{i=0}^{K_{z}} b_{K_{z},i}^{l} E_{t_{n}}^{x} \left[z^{n+i} \right].$$
(11)

Zhao [19] showed that in order to have a stable semi-discrete scheme, the following should hold:

$$k = K_y$$
, with $K_y = 1, 2, ..., 7$ and $K_y = 9$
 $l = 1$, with $K_z = 1, 2, 3$. (12)

The coefficients are presented in the following Table 1 and 2.

2.2 The stable fully discrete scheme

Let $\mathbb{R}_{\Delta x}$ denote a partition of the real axis, i.e. $\mathbb{R}_{\Delta x} = \{x_i | x_i \in \mathbb{R}, i \in \mathbb{Z}, x_i < x_{i+1}, \lim_{i \to +\infty} x_i = +\infty, \lim_{i \to -\infty} x_i = -\infty\}$. The fully discrete scheme is defined as (cf. [19]): Given random variables $(y_i^{N-l}, z_i^{N-l}), l = 0, 1, \dots, K-1$ with K =

Lorenc Kapllani, Long Teng and Matthias Ehrhardt

Ta	ble 1: The	coefficients	$\{b_{K_y,i}^{K_y}\}_{i=0}^{K_y}$ u	ntil <i>K</i> _y
		$b_{K_y,i}^{K_y}$		
i = 0	i = 1	<i>i</i> = 2	<i>i</i> = 3	
$\frac{1}{2}$	$\frac{1}{2}$	1		
$\frac{1}{6}$	4 3 8	$\frac{1}{6}$ $\frac{3}{8}$	$\frac{1}{8}$	
			$b_{K_y,i}^{K_y}$	Table 1: The coefficients $\{b_{K_y,i}^{K_y}\}_{i=0}^{K_y}$ u $b_{K_y,i}^{K_y}$ $i=0$ $i=1$ $i=2$ $i=3$ $\frac{1}{2}$

Table 2: The coefficients $\{b_{K_z,i}^1\}_{i=0}^{K_z}$ for $K_z = 1, 2, 3$

Kz			$b^1_{K_z,i}$		
	i = 0	i = 1	i = 2	<i>i</i> = 3	
1 2 3	$\frac{\frac{1}{2}}{\frac{5}{12}}$	$\frac{\frac{1}{2}}{\frac{8}{12}}$	$-\frac{1}{12}$ $-\frac{5}{24}$	$\frac{1}{24}$	

 $\max\{K_y, K_z\}$, find the random variables (y_i^n, z_i^n) , $n = N - K, \dots, 0$ such that

$$y_{i}^{n} = \hat{E}_{t_{n}}^{x_{i}} \left[\hat{y}^{n+K_{y}} \right] + K_{y} \Delta t \sum_{j=1}^{K_{y}} b_{K_{y},j}^{K_{y}} \hat{E}_{t_{n}}^{x_{i}} \left[f(t_{n+j}, \hat{y}^{n+j}, \hat{z}^{n+j}) \right] + K_{y} \Delta t b_{K_{y},0}^{K_{y}} f(t_{n}, y_{i}^{n}, z_{i}^{n})$$

$$0 = \hat{E}_{t_{n}}^{x_{i}} \left[\hat{z}^{n+1} \right] + \sum_{j=1}^{K_{z}} b_{K_{z},j}^{1} \hat{E}_{t_{n}}^{x_{i}} \left[f(t_{n+j}, \hat{y}^{n+j}, \hat{z}^{n+j}) \Delta W_{t_{n+j}} \right]$$

$$- \sum_{j=1}^{K_{z}} b_{K_{z},j}^{1} \hat{E}_{t_{n}}^{x_{i}} \left[\hat{z}^{n+j} \right] - b_{K_{z},0}^{1} z_{i}^{n},$$
(13)

where (y_i^n, z_i^n) denotes the approximation of $(y(t_n, x_i), z(t_n, x_i))$, $\hat{E}_{t_n}^{x_i}[\cdot]$ is the approximation of $E_{t_n}^{x_i}[\cdot]$ and $(\hat{y}^{n+j}, \hat{z}^{n+j})$ are the interpolating values from (y^{n+j}, z^{n+j}) at the space point $x_i + W_{t_{n+j}} - W_{t_n}$. In order to approximate the conditional expectations, the Gauss-Hermite quadrature rule is used, due to the high accuracy that can be achieved only with a few points. Therefore, the conditional expectation can be expressed as

$$\hat{E}_{t_n}^{x_i}[\hat{y}^{n+k}] = \frac{1}{\sqrt{\pi}} \sum_{j=1}^{L} \omega_j \hat{y}^{n+k} (x_i + \sqrt{2k\Delta t} \, a_j), \tag{14}$$

where (ω_i, a_i) , for j = 1, ..., L are the weights and roots of the Hermite polynomial of degree L (see [19]). In the same way, one can express the other conditional expectations in (13).

The error of the method exhibits different behaviours due to the time-space discretization. However, the maximal order of convergence is 3 for both processes. For technical details, we refer to [19]. Due to the high order of convergence, the numerical method can achieve very high accuracy. It can be observed from (13) that the calculations on each point are independent. Therefore parallelization techniques can be easily adapted. In the next Section 3, we discuss the algorithmic framework.

3 The Algorithmic Framework

3.1 The Algorithm

According to Section 2, the whole process for solving (2) is divided into 3 steps.

1. Construct the time-space discrete domain.

We divide the time period [0, T] into *N* time steps using $\Delta t = T/N$ and get N + 1 time layers. Moreover, in order to balance the errors in time and space directions, we adjust the space step size Δx and the time step size Δt such that they satisfy the equality $(\Delta x)^r = (\Delta t)^{q+1}$, where $q = \min(K_y + 1, K_z)$ and *r* denotes the global error from the interpolation method used to generate the non-grid points when calculating the conditional expectations.

- 2. Calculate *K* initial solutions with $K = \max\{K_y, K_z\}$. Since only the terminal value is given, one needs to generate the other K - 1 values. This can be done by running a 1-step scheme for $[t_{N-K+1}, t_{N-1}]$ with a higher number of time points such that the K - 1 produced initial values will have neglectable error.
- 3. Calculate the numerical solution (y_0, z_0) backward using equation (13). Note that the calculation for the y_t process is done implicitly by Picard iteration.

3.2 Preliminary considerations

In the numerical experiments, we have considered the following points:

- The space domain needs to be truncated. Since the space domain represent the Brownian motion discretization, in our test we use [-16, 16].
- When generating the non-grid points for the calculation of conditional expectations, some will be outside of the domain. For such points, the value on the boundaries is considered, as the desired solution will not be affected.
- Due to uniformity of the grid, one does not need to consider 2K (*K* for y_t and *K* for z_t) interpolations for each new calculation, but only 2. This is due to the following:

Suppose we are at time layer t_{n-K} . To calculate y_t and z_t values on this time layer, one needs the calculation of conditional expectations for *K* time layers. The cubic spline interpolation is used to find the non-grid values, and the necessary linear systems are solved. For instance, the coefficients for y_t process are $A_y \in$

 $\mathbb{R}^{K \times M}$. All the spline coefficients are stored. When we are at time layer t_{n-K-1} , only the spline interpolation corresponding to the previous calculated values is considered. Then, the columns of matrix A_y are shifted +1 to the right in order to delete the last column and enter the current calculated coefficients in the first column. The new A_y is used for the current step. The same procedure is followed until t_0 . This reduces the amount of work for the algorithm.

• There is a very important benefit from the uniformity of the grid. When we need to find the position of the non-grid point, a naïve search algorithm is to loop over the grid points. In the worst case, a O(M) work is needed. However, this can be done in O(1), i.e. the for loop is removed. Recall that each new point is generated as $X_j = x_i + \sqrt{2\Delta tk} a_j$. This means that taking $int((X_j - x_{min})/\Delta x)$ gives the left boundary of the grid interval that X_j belongs to. This reduces the total computation time substantially, as it will be demonstrated in the numerical experiments.

3.3 The Parallel implementation

In this Section we present the naïve parallelization of the multistep scheme. Nevertheless, we have kept into attention the optimal CUDA execution model, i.e. creating arrays such that the access will be aligned and coalesced, reducing the redundant access to global memory, using registers when needed etc.

The first and second steps of the algorithm are implemented in the host. The third step is fully implemented in the device. Recall from (13) that the following steps are needed to calculate the approximated values on each time layer backward:

• Generation of non-grid points $X_i = x_i + \sqrt{2\Delta t k a_i}$.

In the uniform domain, the non-grid points need to be generated only once. To do this, a kernel is created where each thread generates L points.

• Calculation of the values \hat{y} and \hat{z} at the non-grid points.

This is the most time consuming part of the algorithm, since it involves the solution of two linear systems (see third point in Subsection 3.2) arising from the spline interpolation.

We used the BiCGSTAB iterative method since the matrix is tridiagonal. To apply the method, we considered the cuBLAS and cuSPARSE libraries. For the inner product, second norm and addition of vectors, we use the cuBLAS library.

For the matrix vector multiplication, we use the cuSPARSE library with the compressed sparse row format, due to the structure of the system matrix. Moreover, we created a kernel to calculate the spline coefficients based on the solved systems under the spline interpolation idea.

Finally, a kernel to apply the last point in Subsection 3.2 was created to find the values at non-grid points. Note that each thread is assigned to find the values.

• Calculation of the conditional expectations.

For the first conditional expectations in the right hand side of (13), we created one kernel, where each thread calculates one value by using (14). Moreover, we merged the calculation of three conditional expectation in one kernel, namely

 $\hat{E}_{t_n}^{x_i}[\hat{z}^{n+j}], \quad \hat{E}_{t_n}^{x_i}[f(t_{n+j}, \hat{y}^{n+j}, \hat{z}^{n+j})], \quad \hat{E}_{t_n}^{x_i}[f(t_{n+j}, \hat{y}^{n+j}, \hat{z}^{n+j})\Delta W_{t_{n+j}}],$

for j = 1, 2, ..., K. This reduces the accessing of data multiple times from the global memory. Note that one thread calculates three values as in (14).

• Calculation of the *z*^{*t*} values.

The second equation in (13) is used and each thread calculates one value.

• Calculation of the *y*_t values.

The first equation in (13) is used and each thread calculates one value, using the Picard iterative process.

4 Numerical Results

We implement the parallel algorithm using CUDA C programming. The parallel computing times are compared with the serial ones on a CPU. Furthermore, the speedups are calculated. The CPU is Intel(R) Core(TM) i5-4670 3.40Ghz with 4 cores. The GPU is a NVIDIA GeForce 1070 Ti with a total 8GB GDDR5 memory.

In the following we consider an option pricing example, the Black-Scholes model. Consider a security market that contains one bond with price p_t and one stock with price S_t . Therefore, their dynamics are described by:

$$\begin{cases} dp_t = r_t p_t dt, \quad t \ge 0, \\ p_0 = p, \end{cases}$$
(15)

$$\begin{cases} dS_t = \mu_t S_t dt + \sigma_t S_t dW_t, \quad t \ge 0, \\ S_0 = x, \end{cases}$$
(16)

where r_t denotes the interest rate of the bond, p is its current value, μ_t is the expected return on the stock S_t , σ_t is the volatility of the stock, x is its current value and W_t denotes the Brownian motion.

An European Option is a contract that gives the owner the right, but not the obligation, to buy or sell the underlying security at a specific price, known as the strike price K, on the option's expiration date T. A European call option gives the owner the right to purchase the underlying security, while a European put option gives the owner the right to sell the underlying security. Let us take the European call option as an example. The decision of the holder will depend on the stock price at maturity T. If the value of the stock $S_T < K$, then the holder would discard the option; whereas if $S_T > K$, the holder would use the option and make a profit of $S_T - K$. Therefore, the payoff of a call option is $(S_T - K)^+$ and for a put option $(K - S_T)^+$, where $(f)^+ = \max(0, f)$. The option pricing problem of the writer (or seller) is to determine a premium for this contract at present time t_0 . Note that the payoff function is an $\{\mathscr{F}_T\}$ -measurable random variable.

Suppose that an agent sells the option at price y_t and then invests it in the market. Denote his wealth on each time by y_t . Assume that at each time the agent invests a portion of his wealth in an amount given by π_t into the stock, and the rest $(y_t - \pi_t)$ into the bond. Now the agent has a portfolio based on the stock and the bond. Considering a stock that pays a dividend $\delta(t, S_t)$, the dynamics of the wealth process y_t are described by

$$dy_{t} = \frac{\pi_{t}}{S_{t}} dS_{t} + \frac{y_{t} - \pi_{t}}{p_{t}} dp_{t} + \pi_{t} \delta(t, S_{t}) dt$$

$$= \frac{\pi_{t}}{S_{t}} (\mu_{t} S_{t} dt + \sigma_{t} S_{t} dW_{t}) + \frac{y_{t} - \pi_{t}}{p_{t}} (r_{t} p_{t} dt) + \pi_{t} \delta(t, S_{t}) dt$$

$$= (r_{t} y_{t} + \pi_{t} (\mu_{t} - r_{t} + \delta(t, S_{t}))) dt + \pi_{t} \sigma_{t} dW_{t}.$$
(17)

Let $z_t = \pi_t \sigma_t$, then

$$-dy_t = -\left(r_t y_t + \left(\mu_t - r_t + \delta(t, S_t)\right)\frac{z_t}{\sigma_t}\right)dt + z_t \, dW_t.$$
(18)

For a call option, one needs to solve a Forward Backward Stochastic Differential Equation (FBSDE), where the forward part is given from the SDE modelling of the stock price dynamics.

Example 1. Let us consider the Black-Scholes FBSDE

$$\begin{cases} dS_t = \mu_t S_t dt + \sigma_t S_t dW_t, \quad S_0 = x, \quad t \in [0, T] \\ -dy_t = -\left(r_t y_t + \left(\mu_t - r_t + \delta(t, S_t)\right) \frac{z_t}{\sigma_t}\right) dt + z_t dW_t, \quad t \in [0, T) \\ y_T = (S_T - K)^+. \end{cases}$$
(19)

For constant parameters (i.e. $r_t = r$, $\mu_t = \mu$, $\sigma_t = \sigma$, $\delta_t = \delta$), the analytic solution is

$$\begin{cases} y_t = V(t,S_t) = S_t \exp\left(-\delta(T-t)\right) N(d_1) - K \exp\left(-r(T-t)\right) N(d_2), \\ z_t = \frac{\partial V}{\partial S} \sigma = S_t \exp\left(-\delta(T-t)\right) N(d_1) \sigma, \\ d_{1/2} = \frac{\ln\left(\frac{S_t}{K}\right) + \left(r \pm \frac{\sigma^2}{2}\right)(T-t)}{\sigma\sqrt{T-t}}, \end{cases}$$
(20)

where $N(\cdot)$ is the cumulative standard normal distribution function. In this example, we consider T = 0.33, $K = S_0 = 100$, r = 0.03, $\mu = 0.05$, $\delta = 0.04$, $\sigma = 0.2$, with the solution at (t_0, S_0) being $(y_0, z_0) \doteq (4.3671, 10.0950)$.

Note that the terminal condition has a non-smooth problem for the z_t process. Therefore, for discrete points near the strike price (also called at the money region), the initial value for the z_t process will cause large errors on the next time layers. To overcome this non-smoothness problem, we considered smoothing the initial conditions, cf. the approach of Hendricks [10]. For the forward part of (20), we have the analytic solution

$$S_t = S_0 \exp\left(\left(\mu - \frac{\sigma^2}{2}\right)t + \sigma W_t\right).$$
(21)

10

Discretizing (21), the exponential term will lead to a non-uniform grid. Therefore, instead of working in the stock price domain, we work in the log stock price domain. If we denote $X_t = \ln S_t$, then the analytic solution of X_t reads

$$X_t = X_0 + \left(\mu - \frac{\sigma^2}{2}\right)t + \sigma W_t.$$
(22)

The backward part is the same as the (19). In Table 3 we show the importance of using the log stock price. Note that the speedup is relative to the serial case using a for loop.



Fig. 1: Results of naïve parallelization for the Black-Scholes example.

The naïve results using 256 threads per block are presented in Figure 1 (note that $K_y = K_z = K$). It can be easily observed that the higher accuracy can be achieved when considering a 3-step scheme. Since we have more time layers to consider, more work can be assigned to the GPU and therefore increasing the speedup of the application. The highest speedup that we obtained for the Black-Scholes example is $17 \times .$

Table 3: Comparison due to uniformity of the domain under log stock price transformation for the Black-Scholes model for N = 256, $K = K_y = K_z = 3$ and M = 24826.

Туре	Time	Speedup
Serial (with for)	36825.27	
Serial (without for)	76.94	478.62
Parallel (with for)	237.57	155.01
Parallel (without for)	6.72	5476.19

==22320== Profiling result:							
Type	Time(%)	Time	Calls	Avg	Min	Max	Name
GPU activities:	58.95%	7.24543s	85124	85.116us	15.520us	207.81us	void nrm2_kernel <double, dou<="" td=""></double,>
	7.88%	968.66ms	1527	634.36us	595.40us	716.39us	<pre>sp_inter_non_grid_d_no_for(c</pre>
	7.39%	908.40ms	1527	594.89us	520.64us	763.43us	<pre>gen_non_grid(double*, double</pre>
	7.12%	875.58ms	42562	20.571us	19.040us	237.67us	void csrMv_kernel <double, do<="" td=""></double,>
	5.45%	670.20ms	1527	438.90us	388.77us	494.69us	<pre>calc_f_and_c_exp_d(double*,</pre>
	3.55%	436.25ms	85124	5.1240us	3.5840us	8.8000us	void dot_kernel <double, doub<="" td=""></double,>
	1.62%	199.67ms	42562	4.6910us	3.4560us	8.0000us	void axpy_kernel_val <double,< td=""></double,<>
	1.62%	199.27ms	85124	2.3400us	1.8880us	7.9040us	void reduce_1Block_kernel <dc< td=""></dc<>
	1.29%	159.04ms	509	312.45us	308.07us	333.60us	<pre>calc_c_exp_d(double*, double</pre>
	1.19%	146.80ms	21569	6.8060us	4.4160us	8.7360us	step_3(double*, double*, dou
	1.03%	126.59ms	42571	2.9730us	544ns	34.647ms	[CUDA memcpy HtoD]
	0.68%	83.450ms	127687	653ns	608ns	7.6160us	[CUDA memcpy DtoH]
	0.67%	82.697ms	21569	3.8340us	3.2960us	7.7440us	step_8(double*, double*, dou
	0.42%	51.357ms	12192	4.2120us	3.4560us	6.2720us	copy_d(double*, double*, int
	0.39%	48.323ms	509	94.936us	91.937us	103.94us	<pre>calc_y(double*, double*, double*, double*</pre>
	0.37%	45.056ms	20993	2.1460us	1.8880us	5.3760us	step_13(double*, double*, do
	0.12%	15.323ms	6144	2.4930us	1.5360us	5.8880us	void copy_kernel <double, int<="" td=""></double,>
	0.12%	14.548ms	512	28.414us	26.560us	35.040us	<pre>spline_coeff(double*, double</pre>
	0.07%	8.9152ms	509	17.515us	16.512us	19.680us	<pre>calc_z(double*, double*, dou</pre>
ouble*, int, doub	le)						
	0.00%	15.168us	1	15.168us	15.168us	15.168us	CSR(double*, int*, int*, int
	0.00%	4.0640us	1	4.0640us	4.0640us	4.0640us	<pre>csinterp_d(double*, double*,</pre>
	0.00%	736ns	1	736ns	736ns	736ns	[CUDA memset]

(a) Performance of naïve kernels

==22383== Profiling result:							
Туре	Time(%)	Time	Calls	Avg	Min	Max	Name
GPU activities:	18.57%	960.11ms	1527	628.75us	588.39us	899.21us	<pre>sp_inter_non_grid_d_no_for(</pre>
	17.69%	914.31ms	1527	598.76us	521.16us	765.06us	<pre>gen_non_grid(double*, doubl</pre>
	16.82%	869.32ms	42562	20.424us	18.880us	23.584us	void csrMv_kernel <double, d<="" td=""></double,>
	12.76%	659.30ms	1527	431.76us	387.52us	487.91us	calc_f_and_c_exp_d(double*,
	11.24%	580.94ms	127686	4.5490us	3.4240us	9.1200us	void dot_kernel <double, dou<="" td=""></double,>
	5.27%	272.61ms	127686	2.1340us	1.8560us	114.79us	<pre>void reduce_1Block_kernel<d< pre=""></d<></pre>
	3.42%	176.80ms	42562	4.1530us	2.9440us	7.7760us	<pre>void axpy_kernel_val<double< pre=""></double<></pre>
	3.08%	159.16ms	509	312.69us	308.83us	334.11us	<pre>calc_c_exp_d(double*, doubl</pre>
	2.72%	140.38ms	21569	6.5080us	4.4480us	8.0320us	<pre>step_3(double*, double*, do</pre>
	1.86%	95.952ms	9	10.661ms	544ns	34.389ms	[CUDA memcpy HtoD]
	1.53%	79.093ms	127687	619ns	576ns	7.5840us	[CUDA memcpy DtoH]
	1.50%	77.376ms	21569	3.5870us	3.1360us	7.4240us	<pre>step_8(double*, double*, do</pre>
	0.97%	50.224ms	12192	4.1190us	3.3600us	7.0080us	copy_d(double*, double*, in
	0.93%	47.838ms	509	93.984us	87.809us	112.77us	calc_y(double*, double*, do
	0.81%	41.774ms	20993	1.9890us	1.7920us	7.2960us	<pre>step_13(double*, double*, d</pre>
	0.29%	15.097ms	6144	2.4570us	1.5040us	7.1040us	void copy_kernel <double, in<="" td=""></double,>
	0.28%	14.292ms	512	27.914us	25.856us	35.840us	<pre>spline_coeff(double*, doubl</pre>
	0.17%	8.7412ms	509	17.173us	16.352us	19.136us	<pre>calc_z(double*, double*, do</pre>
	0.11%	5.5165ms	1024	5.3870us	4.6720us	6.1760us	<pre>rhs(double*, double*, int,</pre>
	0.00%	14.048us	1	14.048us	14.048us	14.048us	CSR(double*, int*, int*, in
	0.00%	3.9040us	1	3.9040us	3.9040us	3.9040us	<pre>csinterp_d(double*, double*</pre>
	0.00%	736ns	1	736ns	736ns	736ns	[CUDA memset]

(b) Performance after first optimization iteration

==22446== Profiling result:							
Type	Time(%)	Time	Calls	Avg	Min	Max	Name
GPU activities:	24.09%	910.16ms	1527	596.05us	552.26us	777.51us	<pre>sp_inter_non_grid_d_no_for(d</pre>
	22.59%	853.73ms	42562	20.058us	18.624us	231.43us	void csrMv_kernel <double, do<="" td=""></double,>
	16.13%	609.41ms	127686	4.7720us	3.5840us	9.0240us	void dot_kernel <double, doub<="" td=""></double,>
	7.92%	299.37ms	127686	2.3440us	1.8560us	8.0320us	<pre>void reduce_1Block_kernel<dc< pre=""></dc<></pre>
	5.90%	222.81ms	1527	145.91us	133.38us	167.65us	<pre>calc_f_and_c_exp_d(double*,</pre>
	4.44%	167.65ms	42562	3.9380us	2.9440us	7.9680us	void axpy_kernel_val <double,< td=""></double,<>
	3.74%	141.18ms	21569	6.5450us	4.5120us	7.6800us	<pre>step_3(double*, double*, dou</pre>
	2.75%	103.89ms	1527	68.035us	61.664us	92.384us	<pre>gen_non_grid(double*, double</pre>
	2.55%	96.399ms	9	10.711ms	544ns	34.512ms	[CUDA memcpy HtoD]
	2.10%	79.491ms	127687	622ns	576ns	7.5520us	[CUDA memcpy DtoH]
	2.06%	77.984ms	21569	3.6150us	3.1680us	6.9120us	<pre>step_8(double*, double*, dou</pre>
	1.29%	48.910ms	12192	4.0110us	2.7200us	7.2320us	copy_d(double*, double*, int
	1.27%	47.908ms	509	94.122us	87.840us	108.10us	<pre>calc_y(double*, double*, dou</pre>
	1.10%	41.495ms	20993	1.9760us	1.7600us	6.8160us	<pre>step_13(double*, double*, do</pre>
	0.90%	34.072ms	509	66.938us	65.377us	73.441us	<pre>calc_c_exp_d(double*, double</pre>
	0.40%	15.041ms	6144	2.4480us	1.5040us	7.8400us	<pre>void copy_kernel<double, int<="" pre=""></double,></pre>
	0.38%	14.304ms	512	27.937us	26.016us	35.040us	spline_coeff(double*, double
	0.24%	8.9084ms	509	17.501us	16.768us	20.032us	<pre>calc_z(double*, double*, dou</pre>
	0.16%	6.0294ms	1024	5.8880us	4.6400us	6.7840us	<pre>rhs(double*, double*, int, d</pre>
	0.00%	13.856us	1	13.856us	13.856us	13.856us	CSR(double*, int*, int*, int
	0.00%	4.0000us	1	4.0000us	4.0000us	4.0000us	csinterp_d(double*, double*,
	0.00%	704ns	1	704ns	704ns	704ns	[CUDA memset]

(c) Performance after second optimization iteration

Fig. 2: Results of iterative parallelization for the Black-Scholes example.

==

Furthermore, we optimize the kernels created for the Black-Scholes BSDE for N = 512, $K_y = 3$ and $K_z = 3$. For this, we used the NVIDIA profiling tools (nvprof and nvvp) to gather information about performance bottlenecks and apply the proper optimization technique.

After applying nvprof, the main bottleneck of the application is the second norm kernel that calculates the errors in the BiCGSTAB algorithm as presented in Figure 2a. Note that this kernel is already optimized by the NVIDIA developers. However, it is a kernel that serves for a general purpose. Instead of using second norm kernel, we used the dot kernel and later took the square root. This reduced the computation time from 7.2 s to 580.9 ms as presented in Figure 2b. We consider it as the first iteration of the optimization process.

Moreover, we applied again nvprof and found that the next performance inhibitor is the kernel which calculates the non-grid values and another kernel that generates the non-grid values. This is due to the inefficient memory accesses. To overcome this problem, we considered loop interchanging and loop unrolling. As presented in Figure 2c, the performance of above kernels is improved, from 960.1 *ms* to 910.1 *ms* and 914.3 *ms* to 103.9 *ms* respectively. Finally, we changed the thread configuration to 128 threads per block in order to increase parallelizm and we were able to achieve a $51 \times$ speedup. We present the speedups for each iteration of the optimization process in Table 4.

Table 4: Speedups of Black-Scholes model for N = 256, $K = K_y = K_z = 3$ and M = 24826.

Туре	Time	Speedup	
Serial	307.09		
Naïve	18.2	16.87	
First iteration	7.45	41.22	
Second iteration	6.02	51.01	

5 Conclusions and Outlook

In this work we parallelized the multistep method developed in [19] for the numerical approximation of BSDEs on GPU.

Firstly, we presented an optimal operation to find the location of the interpolated values. This was essential for the reduction of the computational time. The numerical results exhibited a high accuracy in very small computation times. Moreover, we optimized the application after finding the performance bottlenecks and applying optimization techniques. Using the *cuBLAS* kernel to calculate the error of the BiCGSTAB iterative method, loop interchange and loop unrolling provided us a $51 \times$ speedup for the Black-Scholes example.

Based on our results, the GPU architecture for the multistep scheme is well suited for the acceleration of BSDEs. For future work we will focus on parallelizing *d*dimensional problems using more time steps as in [17] for a higher accuracy with financial applications such as multi-asset option pricing.

Acknowledgements The authors were partially supported by the bilateral German-Portuguese Project *FRACTAL – FRActional models and CompuTationAL Finance*, the bilateral German-Hungarian Project *CSITI – Coupled Systems and Innovative Time Integrators* and the bilateral German-Slovakian Project *ENANEFA – Efficient Numerical Approximation of Nonlinear Equations in Financial Applications* all financed by the DAAD.

References

- Behforooz, G. and Papamichael, N., 1979. End conditions for cubic spline interpolation, J. Inst. Math. Appl., 23, pp. 355-366, 1979.
- Behforooz, G., 1995. A Comparison of the E(3) and Not-A-Knot Cubic Splines, J. Inst. Math. Appl. Math. Comp., 72, pp. 219-223.
- Bender, C. and Steiner, J., 2012. Least-squares Monte-Carlo for backward SDEs. In: Numerical methods in finance (pp. 257-289). Springer, Berlin, Heidelberg.
- Bouchard, B. and Touzi, N., 2004. Discrete-time approximation and Monte-Carlo simulation of backward stochastic differential equations. Stochastic Processes and their applications, 111(2), pp. 175-206.
- Crisan, D. and Manolarakis, K., 2012. Solving backward stochastic differential equations using the cubature method: application to nonlinear pricing. SIAM J. Fin. Math., 3(1), pp.534-571.
- Dai, B., Peng, Y. and Gong, B., 2010. Parallel option pricing with BSDE method on GPU. In 2010 Ninth Int. Conference on Grid and Cloud Computing (pp. 191-195). IEEE.
- El Karoui, N., Peng, S. and Quenez, M.C., 1997. Backward stochastic differential equations in finance. Mathematical Finance, 7(1), pp. 1-71.
- Gobet, E., Lemor, J.P. and Warin, X., 2005. A regression-based Monte Carlo method to solve backward stochastic differential equations. Annal. Appl. Probab., 15(3), pp. 2172-2202.
- Gobet, E., López-Salas, J.G., Turkedjiev, P. and Vázquez, C., 2016. Stratified regression Monte-Carlo scheme for semilinear PDEs and BSDEs with large scale parallelization on GPUs. SIAM J. Sci. Comput., 38(6), pp. C652-C677.
- Hendricks, C., 2017. High-Order Methods for Parabolic Equations in Multiple Space Dimensions for Option Pricing Problems, Dissertation, University of Wuppertal.
- Lemor, J.P., Gobet, E. and Warin, X., 2006. Rate of convergence of an empirical regression method for solving generalized backward stochastic differential equations. Bernoulli, 12(5), pp. 889-916.
- Pardoux, E. and Peng, S., 1990. Adapted solution of a backward stochastic differential equation. Systems & Control Letters, 14(1), pp. 55-61.
- Peng, S., 1991. Probabilistic interpretation for systems of quasilinear parabolic partial differential equations. Stochastics and stochastics reports, 37(1-2), pp. 61-74.
- Peng, Y., Gong, B., Liu, H. and Dai, B., 2011. Option pricing on the GPU with backward stochastic differential equation. In 2011 Fourth International Symposium on Parallel Architectures, Algorithms and Programming (pp. 19-23). IEEE.
- Peng, Y., Liu, H., Yang, S. and Gong, B., 2014. Parallel algorithm for BSDEs based high dimensional American option pricing on the GPU. J. Comp. Inform. Syst., 10(2), pp. 763-771.

14

A Multistep Scheme to solve BSDEs for Option Pricing on GPUs

- Teng, L., 2018. A Review of tree-based approaches to solve forward-backward stochastic differential equations. arXiv preprint arXiv:1809.00325.
- Teng, L., Lapitckii, A. and Gnther, M., 2018. A Multi-step Scheme based on Cubic Spline for solving Backward Stochastic Differential Equations. arXiv preprint arXiv:1809.00324.
- van der Vorst, H.A., 1992. Bi-CGSTAB: A Fast and Smoothly Converging Variant of Bi-CG for the Solution of Nonsymmetric Linear Systems, SIAM J. Sci. Stat. Comput., 13(2), pp. 631-644.
- Zhao, W., Zhang, G. and Ju, L., 2010. A stable multistep scheme for solving backward stochastic differential equations. SIAM J. Numer. Anal., 48(4), pp. 1369-1394.
- Zhao, W., Chen, L. and Peng, S., 2006. A new kind of accurate numerical method for backward stochastic differential equations. SIAM J. Sci. Comput., 28(4), pp. 1563-1581.